

Lecture 2:

If X scheme /S

Stack : "functor" $\text{Sch} \rightarrow \text{Gpd}$.

Moduli stack of vector bundles of rank r on X

$U \in \text{Sch}/S \mapsto$ vector bundles of rank r on $X \times U$, their isomorphisms

If $f: U \rightarrow U'$

$f^*: \text{Bun}_X^r(U) \rightarrow \text{Bun}_X^r(U')$ pull-back via $(\text{id} \times f)$

$$(g \circ f)^* \cong f^* \circ g^*$$

Quotient stack: $G \backslash X$

$[X/G]: U \mapsto G\text{-principal bundles } P \text{ on } U$
 + G -equiv. maps $P \rightarrow X$
 + their isomorphisms

If $f: U \rightarrow U'$, $[X/G](U) \rightarrow [X/G](U')$ pull-back of principal bundle.

I. Grothendieck topology

\mathcal{C} : category with fiber products

A Grothendieck topology on \mathcal{C} is: For each $U \in \mathcal{C}$ a set of families

$\tau(U)$: set of coverings. $\{U_i \rightarrow U\}$ st

1) If $U \rightarrow U$ isomorphism, then $\{U \rightarrow U\} \in \tau(U)$

2) (Transitivity): If $\{U_i \xrightarrow{\phi_i} U\} \in \tau(U)$ and $\forall i: \{U_{ij} \xrightarrow{\psi_{ij}} U_i\} \in \tau(U_i)$

then $\{U_{ij} \xrightarrow{\psi_{ij}} U\}_{i,j} \in \tau(U)$

3) (Base change) If $\{U_i \rightarrow U\} \in \tau(U)$ and $V \rightarrow U$

then $\{V \times_U U_i \rightarrow V\} \in \tau(V)$.

\mathcal{C} + Grothendieck top: Date.

Basic example: X topological space.

$\mathcal{O}(X)$: category of open sets
maps: $V \hookrightarrow U$.

$\tau(U)$ = set of open covers of U
covering = set of families $\{U_i \xrightarrow{\varphi_i} U\}$ st $\bigcup \varphi_i(U_i) = U$.

If X scheme/S we define:

Small Zariski site: Objects $U \hookrightarrow X$ Sdr/S
Morphisms $U \rightarrow U'$
 \sim_{X^2}
covering of U : $\{U_i \xrightarrow{\varphi_i} U\}$ st $\bigcup \varphi_i(U_i) = U$.

Small étale site: Objects: étale morphisms $U \rightarrow X$ defining $U \in \text{Sch}/S$
morphisms: $U \rightarrow U'$ étale morphism +
 \sim_X
étale covering (covers of U : $\{U_i \xrightarrow{\varphi_i} U\}$) φ_i : étale st $\bigcup \varphi_i(U_i) = U$.

Same construction with étale replaced by:

- smooth
- fppf
- fpqc

Remark:

Open embedding \Rightarrow étale \Rightarrow smooth \Rightarrow fppf \Rightarrow fpqc

Big sites • étale: Objects: Sch/S

Morphisms: morphisms of schemes

Covers of X : same.

Def: If \mathcal{C} is a site.

A presheaf on \mathcal{C} is $F: \mathcal{C}^{op} \rightarrow \text{Set}$

A sheaf on \mathcal{C} is a presheaf $F: \mathcal{C}^{op} \rightarrow \text{Set}$ such that

$\forall U \in \mathcal{C}, f, g \in F(U)$ if $\{U_i \xrightarrow{\varphi_i} U\} \in \tau(U)$

$f|_{U_i} = g|_{U_i}$ (here: $F(\varphi_i)(f) = F(\varphi_i)(g)$) $\Rightarrow f = g$.

2. $\forall U \in \mathcal{E}, \forall \{U_i \xrightarrow{\varphi_i} U\} \in \tau(U)$

$\forall f_i \in F(U_i) \text{ st } F(\varphi_{i,j,i})(f_i) = F(\varphi_{i,j,j})(f_j) \text{ in } F(U_i \times_U U_j)$

$[\varphi_{i,j,i}: U_i \times_U U_j \rightarrow U_i]$

" $U_i \times_U U_j \subseteq U_i$ " " " $U_i \times_U U_j \subseteq U_j$ "

$\Rightarrow \exists f \in F(U) \text{ st } \forall i \quad f|_{U_i} = f_i.$

Remark: Sheaf for the fpqc top \Rightarrow sheaf for fppf \Rightarrow smooth \Rightarrow etale \Rightarrow Zar.

If $X \in \text{Sch}/S$ we can define $\text{Hom}(-, X): (\text{Sch}/S)^{\text{op}} \rightarrow \text{Set}$

Theorem: This is a sheaf for the Zariski topology. (tautological)

Theorem (Grothendieck): This is a sheaf for the fpqc topology ("descent theory").

Remark: Schemes \subseteq Alg. spaces \subseteq Stacks \subseteq Functions $\text{Sch}/S^{\text{op}} \rightarrow \text{Set}$.

II. 2-categories

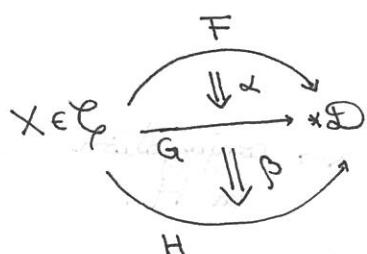
Cat = category of all categories

Objects: categories

Morphisms: functors

+ between two functors $\mathcal{C} \xrightarrow{F} \mathcal{D}$ natural transformations

+ vertical composition of natural transformations



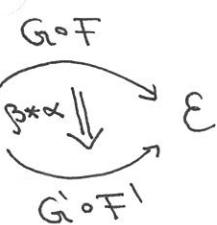
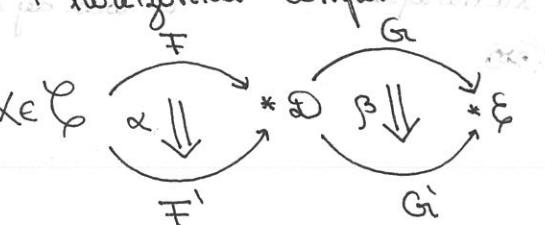
$$\alpha_x: F(x) \rightarrow G(x)$$

$$\beta_x: G(x) \rightarrow H(x)$$

$$\beta \circ \alpha: F \Rightarrow H$$

$$(\beta \circ \alpha)_x: \beta_x \circ \alpha_x$$

+ horizontal composition



$$\beta * \alpha: G \circ F \Rightarrow G' \circ F'$$

$$x \in \mathcal{C}: \alpha_x: F(x) \rightarrow F'(x)$$

$$G \alpha_x: G(F(x)) \rightarrow G(F'(x))$$

$$\beta_{F'(x)}: G(F'(x)) \rightarrow G'(F'(x))$$

$$\beta_{F'(x)} \circ G \alpha_x: G(F(x)) \rightarrow G'(F'(x))$$

Déf: A 2-category \mathcal{C} is:

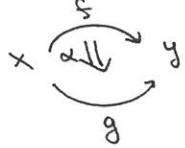
\mathcal{C}_0 : Set of objects = Set of 0-morphisms

\mathcal{C}_1 : Set of 1-morphisms

\mathcal{C}_2 : Set of 2-morphisms

st $\mathcal{C}_0, \mathcal{C}_1$ categories.

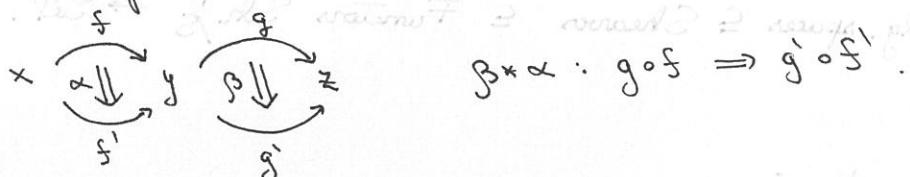
+ between two parallel 1-morphisms $x \xrightarrow{\alpha} y$ there are 2-morphisms



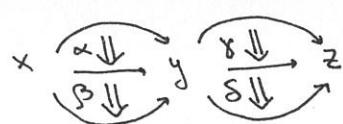
+ identity 2-morphism

+ vertical composition of 2-morphisms, associativity, identity (stabilizing)

+ horizontal composition



+ Axiom:



$$(\gamma \circ \alpha) * (\delta \circ \beta) = (\gamma * \delta) \circ (\alpha * \beta)$$

Example:

• Cat

• Gpd: 1-morphisms are isomorphisms

• Grpd (groups): Here: 2-morphism between $f, g: G \rightarrow H \leftrightarrow$ conjugation in H .

• \mathcal{C} : \mathcal{C}_0 = topological spaces

\mathcal{C}_1 = continuous maps. $f: X \rightarrow Y$

\mathcal{C}_2 = homotopy $\begin{array}{ccc} X & \xrightarrow{\alpha} & Y \\ \Downarrow h & & \end{array}$ modulo homotopies of (\rightarrow vertical composition) or well def.

$H: X \times [0,1] \rightarrow Y$

Other point of view:

A 2-category \Leftrightarrow category \mathcal{C} st $\text{Hom}(x,y)$ is a category (with vertical comp. of 2-morph.)

$$\begin{aligned} \text{Hom}(y,z) \times \text{Hom}(x,y) &\rightarrow \text{Hom}(x,z) \\ (g, f) &\mapsto g \circ f \end{aligned}$$

is a bijunctor (\approx distributivity \Leftrightarrow horizontal composition)

Functor between 2-categories : 2-functors.

All this is called strict 2-categories. (loop to a popular form of 2-categories)

Weak versions:

- Motivation:
- A 1-morphism $f: x \rightarrow y$ is called invertible if $\exists g: y \rightarrow x$ st $gf = \text{id}_x$, $fg = \text{id}_y$. (D o W : 3 laws (W) H are violated)
 - Same for invertible 2-morphism.

→ A 1-morphism $f: x \rightarrow y$ is called an equivalence if $\exists g: y \rightarrow x$ + invertible 2-morphisms $gf \xrightarrow{\sim} \text{id}_x$, $fg \xrightarrow{\sim} \text{id}_y$.

Example: Cat: Functor being an equivalence of categories
Top: Homotopy equivalence.

Recall: $F: \mathcal{C} \rightarrow \mathcal{D}$ is an equiv. of categories
 $\Leftrightarrow \forall y \in \mathcal{D} \ \exists x \in \mathcal{C}$ st $F(x) \xrightarrow{\sim} y$
+ $F: \text{Hom}(a, b) \rightarrow \text{Hom}(F(a), F(b))$ is a bijection.

Dy: A pseudo-functor (weak 2-functor) is $F: \mathcal{C} \rightarrow \mathcal{D}$ between 2-categories st:

$\forall x \in \mathcal{C}, \exists \alpha_x$ invertible 2-morphism with $F(\text{id}_x) \xrightarrow{\alpha_x} \text{id}_{F(x)}$
and $\forall f, g$ 1-morphisms $\exists \epsilon_{g,f}$: invertible 2-morphism st:
 $F(g \circ f) \xrightarrow{\epsilon_{g,f}} F(g) \circ F(f)$

+ F strictly associative on 2-morphisms

+ Axioms on α, ϵ .

Weak 2-category: 2-morphisms, associative composition

For 1-morphisms $h \circ (g \circ f) \xrightarrow{\epsilon_{g,f}} (h \circ g) \circ f$.

\uparrow
invertible
2-morphism

III. Stacks

Def: Let \mathcal{C} be any category. A prestack gives a pseudo-functor

$$\mathcal{C}^{\text{op}} \rightarrow \text{Gpd}$$

Let \mathcal{C} be a site, then a stack is a pseudo-functor $\mathfrak{X}: \mathcal{C} \rightarrow \text{Gpd}$

[Notation: $x \in \mathfrak{X}(U)$ and $f: U' \rightarrow U$, $f_* x|_{U'} = \mathfrak{X}(f)(x)$]

• Glueing of objects:

$$\forall U \in \mathcal{C}, \forall \{U_i \rightarrow U\} \in \tau(U), \forall x_i \in \mathfrak{X}(U_i)$$

$$+ \text{morphisms } \varphi_{ij}: x_i|_{U_{ij}} \rightarrow x_j|_{U_{ij}} \text{ st } \varphi_{ij}|_{U_{ijk}} \circ \varphi_{ik}|_{U_{ijk}} = \varphi_{ik}|_{U_{ijk}}$$

$$\Rightarrow \exists x \in \mathfrak{X}(U) \text{ and } \varphi_i: x|_{U_i} \xrightarrow{\sim} x_i \text{ in } \mathfrak{X}(U_i)$$

$$\text{st } \varphi_{ji} \circ \varphi_i|_{U_{ij}} = \varphi_j|_{U_{ij}}$$

• Isom. form a sheaf:

$$\forall x, y \in \mathfrak{X}(U), \varphi_i: x|_{U_i} \rightarrow y|_{U_i} \text{ st } \varphi_i|_{U_{ij}} = \varphi_j|_{U_{ij}}$$

$$\Rightarrow \exists \eta: x \rightarrow y \text{ st } \eta|_{U_i} = \varphi_i$$

$$\text{and } \forall x, y \in \mathfrak{X}(U) \text{ s.t. } \varphi: x \rightarrow y \text{ s.t. } x \rightarrow y$$

$$\text{st } \varphi|_{U_i} = \eta|_{U_i} \Rightarrow \varphi = \eta$$

$U \mapsto \text{Isom}_U(x, y)$ is a sheaf.